

Counting cycles in planar graphs

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Euler's Formula

A graph drawn in the plane without crossing edges is a **PLANE GRAPH**.

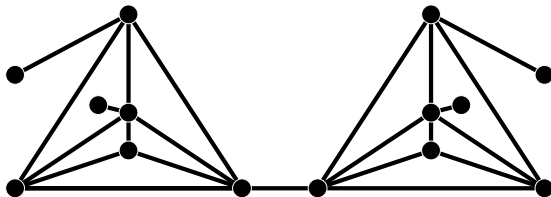
Let G be a plane graph with

- $V = V(G)$, the set of vertices,
- $E = E(G)$, the set of edges,
- $F = F(G)$, the set of faces.

Theorem (Leonhard Euler, 1758)

For any connected plane graph, $|V| - |E| + |F| = 2$.

A graph with a plane graph drawing is a **PLANAR GRAPH**.



Euler's Formula

Theorem (Leonhard Euler, 1758)

For any connected plane graph, $|V| - |E| + |F| = 2$.

Lemma (Handshaking Lemma for plane graphs)

For every plane graph G ,

$$2|E| = \sum_{v \in V} \deg(v) = \sum_{f \in F} \deg(f).$$

Corollary

For every plane graph G with $|V| \geq 3$,

$$|E| \leq 3|V| - 6$$

$$|E| \leq 2|V| - 4, \quad \text{if } G \text{ is bipartite.}$$

Euler's Formula

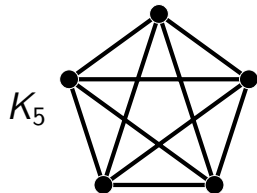
Corollary

For every plane graph G with $|V| \geq 3$,

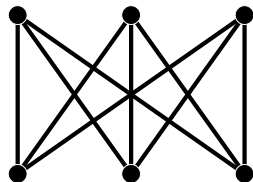
$$|E| \leq 3|V| - 6$$

$$|E| \leq 2|V| - 4,$$

if G is bipartite.



$$|V| = 5, |E| = 10$$



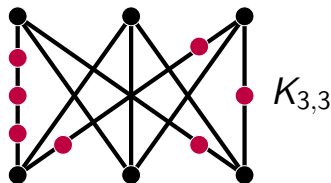
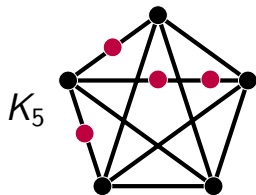
$$|V| = 6, |E| = 9$$

Euler's Formula

Theorem

If G is a graph, then G is planar iff

- G has no **SUBDIVISION** of K_5 or $K_{3,3}$; [Kuratowski, 1930]
- G has no minor of K_5 or $K_{3,3}$. [Wagner, 1937]





Definition

Let $N_{\mathcal{P}}(n, H)$ denote the maximum number of copies of H in an n -vertex PLANAR graph.

Theorem

If $n \geq 3$, then

$$N_{\mathcal{P}}(n, K_2) = 3n - 6.$$

This is achieved by any planar triangulation.

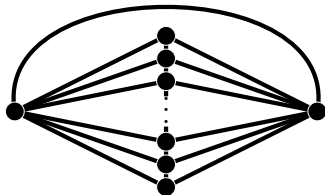
Tutte proved there are $\frac{n^{-7/2}}{64\sqrt{6\pi}} \left(\frac{256}{27}\right)^{n-2}$ planar triangulations.



Theorem (Hakimi-Schmeichel, 1979)

If $n \geq 3$, then

$$N_{\mathcal{P}}(n, C_3) = 3n - 8.$$

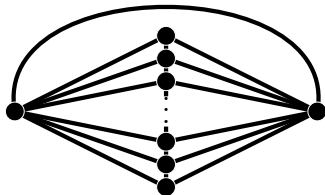


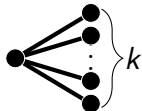


Theorem (Hakimi-Schmeichel, 1979)

If $n \geq 3$, then

$$N_{\mathcal{P}}(n, C_4) = N_{\mathcal{P}}(n, K_{2,2}) = \frac{n^2 + 3n - 22}{2}.$$

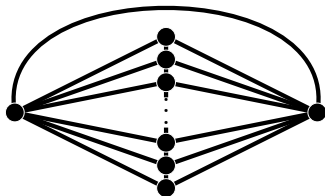


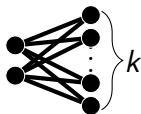


Theorem (Alon-Caro, 1984)

If $k \geq 3$ and $n \geq 4$, then

$$N_{\mathcal{P}}(n, K_{1,k}) = 2 \cdot \binom{n-1}{k} + (n-4) \cdot \binom{4}{k} + 2 \cdot 3k.$$

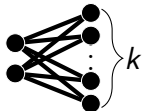




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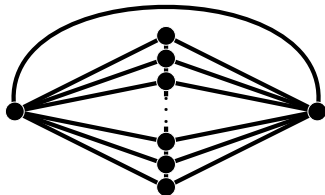
$$N_{\mathcal{P}}(n, K_{2,k}) = \begin{cases} \binom{n-2}{k}, & \text{if } k \geq 5 \text{ or if } k = 4 \text{ and } n \neq 6; \\ 3, & \text{if } k = 4 \text{ and } n = 6; \\ \binom{n-2}{3} + 3(n-4), & \text{if } k = 3 \text{ and } n \neq 6; \\ 12, & \text{if } k = 3 \text{ and } n = 6; \\ \binom{n-2}{2} + 4n - 14, & \text{if } k = 2. \end{cases}$$



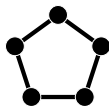
Theorem (Alon-Caro, 1984)

If $n \geq 4$, then

$$N_{\mathcal{P}}(n, K_{2,k}) = \binom{n-2}{k} + O(n^{k-1}).$$

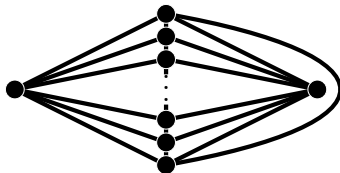


Five-cycles

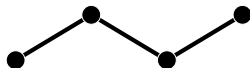


Theorem (Györi-Paulos-Salia-Tompkins-Zamora, 2019)

$$N_{\mathcal{P}}(n, C_5) = \begin{cases} 2n^2 - 10n + 12, & \text{if } n = 6 \text{ or } n \geq 8; \\ 41, & \text{if } n = 7; \\ 6, & \text{if } n = 5. \end{cases}$$



Four-paths



Theorem (Györi-Paulos-Salia-Tompkins-Zamora, 2019)

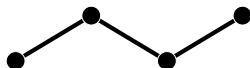
$$N_{\mathcal{P}}(n, P_4) = \begin{cases} 7n^2 - 32n + 27, & \text{if } n \geq 9 \text{ or } n = 5, 6; \\ 222, & \text{if } n = 8; \\ 147, & \text{if } n = 7; \\ 12, & \text{if } n = 4, \end{cases}$$

$$N_{\mathcal{P}}(n, P_k) = \Theta\left(n^{\lfloor (k-1)/2 \rfloor + 1}\right).$$

Theorem (Huynh-Joret-Wood, 2022, generalized by C.-H. Liu, 2021+)

For all H , there exists a ℓ such that $N_{\mathcal{P}}(n, H) = \Theta(n^{\ell})$.

Four-paths



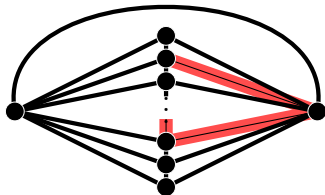
Theorem (Györi-Paulos-Salia-Tompkins-Zamora, 2019)

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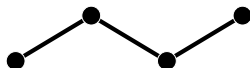
$$N_{\mathcal{P}}(n, P_4) = 7n^2 - 32n + 27.$$

$$7n^2 - 32n + 27$$

$$= 2 \cdot 2 \cdot (n-3) \cdot (n-4) + 2 \cdot (n-2) \cdot (n-3) + 2 \cdot \binom{n-2}{2} + O(n)$$



Four-paths



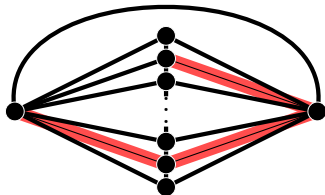
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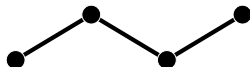
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Four-paths



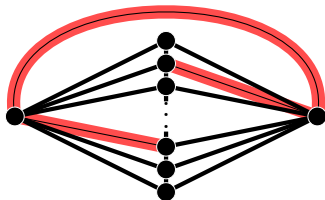
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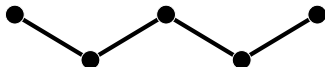
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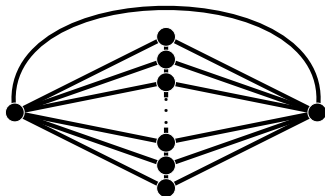


Five-paths



Theorem (Ghosh-Györi-M.-Paulos-Salia-Xiao-Zamora, 2021)

$$N_{\mathcal{P}}(n, P_5) = n^3 + O(n^2)$$



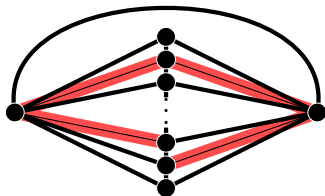
Five-paths

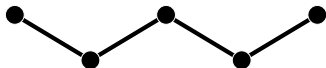


Theorem (Ghosh-Györi-M.-Paulos-Salia-Xiao-Zamora, 2021)

$$N_{\mathcal{P}}(n, P_5) = n^3 + O(n^2)$$

$$n^3 + O(n^2) = (n-2)(n-3)(n-4) + O(n^2)$$





Theorem (Ghosh-Györi-M.-Paulos-Salia-Xiao-Zamora, 2021)

$$N_{\mathcal{P}}(n, P_5) = n^3 + O(n^2)$$

Big idea of the proof:

Lemma

Let $n \geq k \geq 3$ and let G be a planar graph on n vertices such that $S \subseteq V(G)$, $|S| = k$. Then,

$$\sum_{v \in S} \deg(v) \leq 2n + 6k - 16.$$

Proof of five-paths

Theorem (Ghosh-Györi-M.-Paulos-Salia-Xiao-Zamora, 2021)

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Let $n \geq k \geq 3$ and let G be a planar graph on n vertices such that $S \subseteq V(G)$, $|S| = k$. Then,

$$\sum_{v \in S} \deg(v) \leq 2n + 6k - 16.$$

Proof.

$$\sum_{v \in S} \deg(v) = \sum_{v \in S} \deg_{G[S]}(v) + e(S, V - S) \leq 2(3k - 6) + (2n - 4).$$



Proof of five-paths

Theorem (Ghosh-Györi-M.-Paulos-Salia-Xiao-Zamora, 2021)

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Lemma

Let $n \geq k \geq 3$ and let G be a planar graph on n vertices such that $S \subseteq V(G)$, $|S| = k$. Then,

$$\sum_{v \in S} \deg(v) \leq 2n + 6k - 16.$$

If the vertices of G are $\deg(v_1) \geq \deg(v_2) \geq \dots \geq \deg(v_n)$, then

$$\#P_5(G) \leq \sum_{i < j} \deg(v_i) \deg(v_i, v_j) \deg(v_j)$$

Theorem (Ghosh-Györi-M.-Paulos-Salia-Xiao-Zamora, 2021)

$$N_{\mathcal{P}}(n, P_5) = n^3 + O(n^2)$$

If the vertices of G are $\deg(v_1) \geq \deg(v_2) \geq \dots \geq \deg(v_n)$, then

$$\begin{aligned} \#P_5(G) &\leq \sum_{i < j} \deg(v_i) \deg(v_i, v_j) \deg(v_j) \\ &\leq \sum_{i < j} \deg(v_i) \min\{\deg(v_i), \deg(v_j)\} \deg(v_j) \\ &\leq \sum_{i < j} \deg(v_i) (\deg(v_j))^2 \end{aligned}$$

Theorem (Ghosh-Györi-M.-Paulos-Salia-Xiao-Zamora, 2021)

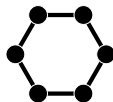
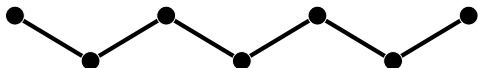
$$N_{\mathcal{P}}(n, P_5) = n^3 + O(n^2)$$

$$\left\{ \begin{array}{l} \text{maximize} \\ \text{subject to} \end{array} \right. \quad \begin{array}{l} \sum_{i < j} x_i x_j^2 \\ n \geq x_1 \geq x_2 \geq \dots \geq x_n \geq 0. \\ \sum_{i=1}^k x_i \leq 2n + 6k - 16, \quad \forall k \geq 3 \\ \sum_{i=1}^n x_i \leq 6n - 12. \end{array}$$

The solution is $n^3 + O(n^2)$.

(Our proof requires $n \geq 11664$.)

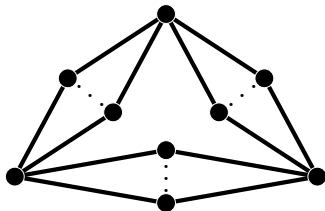
Seven-paths and six-cycles



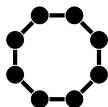
Theorem (Cox-M., 2022)

$$N_{\mathcal{P}}(n, P_7) = \frac{4}{27}n^4 + O(n^{4-1/5})$$

$$N_{\mathcal{P}}(n, C_6) = \left(\frac{n}{3}\right)^3 + O(n^{3-1/5})$$

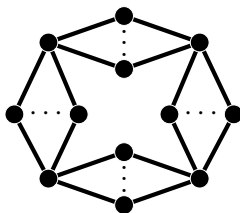


Eight-cycles



Theorem (Cox-M., 2022)

$$N_{\mathcal{P}}(n, C_8) = \left(\frac{n}{4}\right)^4 + O\left(n^{4-1/5}\right)$$

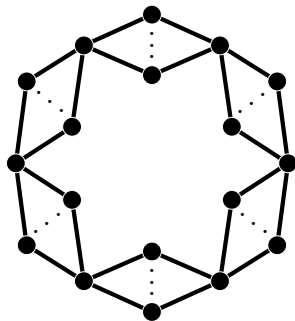
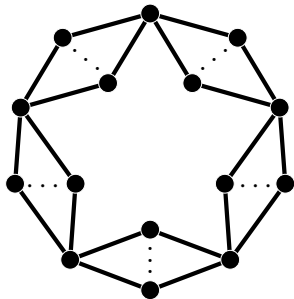


10- and 12-cycles

Theorem (Cox-M., 2023)

$$N_{\mathcal{P}}(n, C_{10}) = \left(\frac{n}{5}\right)^5 + o(n^5)$$

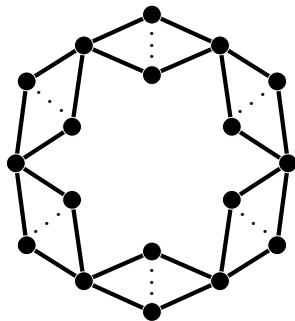
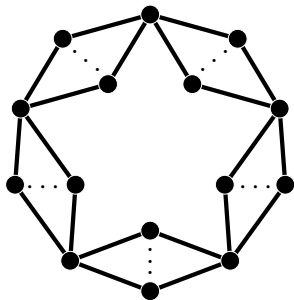
$$N_{\mathcal{P}}(n, C_{12}) = \left(\frac{n}{6}\right)^6 + o(n^6).$$



10- and 12-cycles

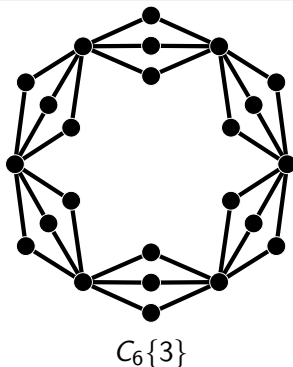
Theorem (Lv-Györi-He-Salia-Tompkins-Zhu, 2024)

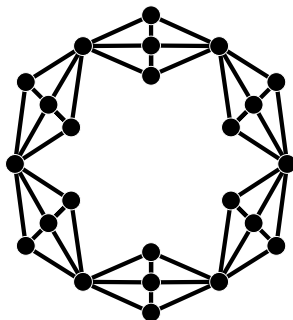
$$N_{\mathcal{P}}(n, C_{2k}) = \binom{n}{k}^k + o(n^k).$$



Definition

Let H be a planar graph. $H\{k\}$ denotes the “planar blow-up of H by k ”. That is, replace each edge of H with a $K_{2,k}$.





Objective

A “reduction lemma” that says $N_{\mathcal{P}}(n, H)$ is asymptotically maximized by

- taking some graph G ,
- blowing up the edges by different amounts, and
- putting a path inside each blowup.

Definition

- Let V be vertex set and let $\mu : \binom{V}{2} \rightarrow [0, 1]$ be a probability mass. That is, $\sum \left\{ \mu(e) : e \in \binom{V}{2} \right\} = 1$.

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Probability mass on a graph

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Theorem (Even Cycle Reduction Lemma)

Let $m \geq 3$.

- $\beta(\mu; m) \stackrel{\text{def}}{=} \sum_{C \in \mathcal{C}(K_V)} \prod_{e \in E(C)} \mu(e) = \mathbb{P}(m \text{ edges form } C_m)$
- $\beta(m) \stackrel{\text{def}}{=} \sup \left\{ \beta(\mu; m) : \text{supp } \mu \subset \binom{V}{2} \text{ for a finite } V \right\}$
- $N_{\mathcal{P}}(n, C_{2m}) \leq \beta(m) \cdot n^m + O(n^{m-1/5})$

Probability mass on a graph

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- Let V be vertex set and let $\mu : \binom{V}{2} \rightarrow [0, 1]$ be a probability mass.
- For a graph $G = (V, E)$, define $\mu(G) = \prod_{e \in E} \mu(e)$.
- For a vertex $x \in V$, let $\bar{\mu}(x) = \sum_{y \in V - \{x\}} \mu(xy)$.

Theorem (Odd Path Reduction Lemma)

Let $m \geq 2$. Then $N_{\mathcal{P}}(n, P_{2m+1}) \leq \rho(m) \cdot n^{m+1} + O(n^{m+4/5})$ where

$$\rho(\mu; m) \stackrel{\text{def}}{=} \frac{1}{2} \sum_{v \in (V)_m} \bar{\mu}(x_1) \left(\prod_{i=1}^{m-1} \mu(x_i x_{i+1}) \right) \bar{\mu}(x_m)$$

Notes about reduction lemmas

- Reduction lemmas are not structural, but counting lemmas.
- We can't guarantee that, e.g., μ isn't positive on a K_5 , non-planar.
- Karush-Kuhn-Tucker (Lagrange multipliers) used to compute ρ or β .

Notes about reduction lemmas

- Reduction lemmas are not structural, but counting lemmas.
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Examples for the cycle C_{2m}

- If μ is an equal distribution on $E(K_m)$, then

$$\sum_{C \in \mathcal{C}(K_m)} \prod_{e \in E(C)} \mu(e) = \frac{(m-1)!}{2} \left(\frac{1}{\binom{m}{2}} \right)^m.$$

- If μ is an equal distribution on $E(C_m)$, then

$$\sum_{C \in \mathcal{C}(K_m)} \prod_{e \in E(C)} \mu(e) = \left(\frac{1}{m} \right)^m \geq \frac{(m-1)!}{2} \left(\frac{1}{\binom{m}{2}} \right)^m$$

Examples for the cycle C_{2m}

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The largest is $\left(\frac{1}{m} \right)^m$ for $m \geq 3$.

- $m = 3, 4, 5, 6$ [Cox-M.]
- $m \geq 7$ [Lv-Györi-He-Salia-Tompkins-Zhu].

Odd cycle Reduction Lemma

- Let V be vertex set and let $\mu : \binom{V}{2} \rightarrow [0, 1]$ be a probability mass.
- For a graph $G = (V, E)$, define $\mu(G) = \prod_{e \in E} \mu(e)$.
- For a vertex $x \in V$, let $\bar{\mu}(x) = \sum_{y \in V - \{x\}} \mu(xy)$.

Theorem (Odd Cycle Reduction Lemma: Heath-M-Wells, 2023+)

Let $m \geq 2$.

- $$\beta(\mu; m) \stackrel{\text{def}}{=} 2m \cdot \sum_{C \in \mathcal{C}(K_V)} \prod_{e \in E(C)} \mu(e) + \sum_{P \in \mathcal{P}(K_V)} \prod_{e \in E(P)} \mu(e)$$
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Theorem (Odd Cycle Reduction Lemma: Heath-M-Wells, 2023+)

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Theorem (Heath-M.-Wells, 2023+)

$$\beta(m) = \frac{2}{m^{m-1}}, \quad \text{if } m \in \{3, 4\};$$
$$\beta(m) \leq \frac{2.7}{m^{m-1}}, \quad \text{if } m \geq 5.$$

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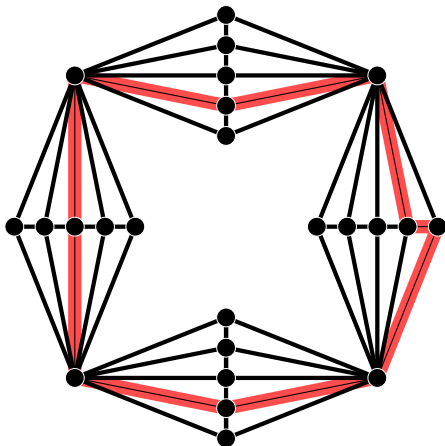
$$N_{\mathcal{P}}(n, C_{2m+1}) = 2m \cdot \left(\frac{n}{m}\right)^m + O\left(n^{m-1/5}\right), \quad \text{if } m \in \{3, 4\};$$

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Odd cycle Reduction Lemma

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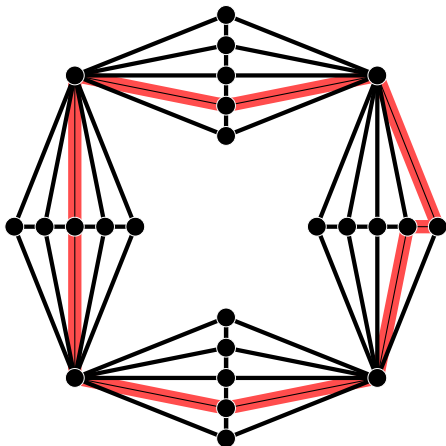
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Odd cycle Reduction Lemma

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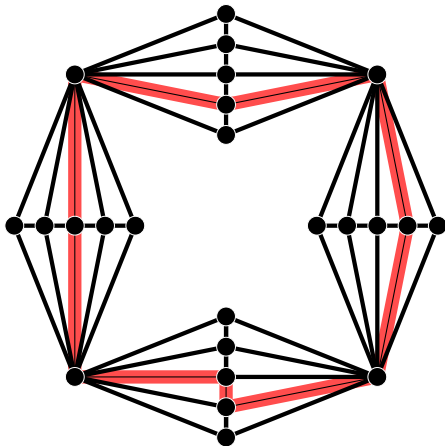
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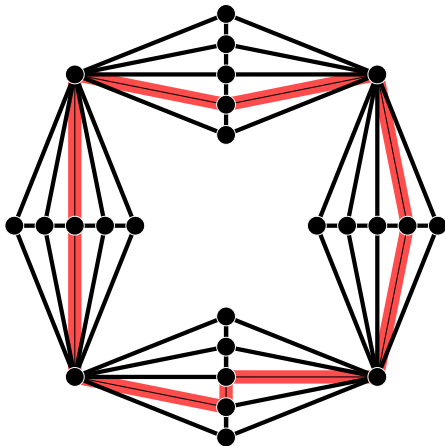
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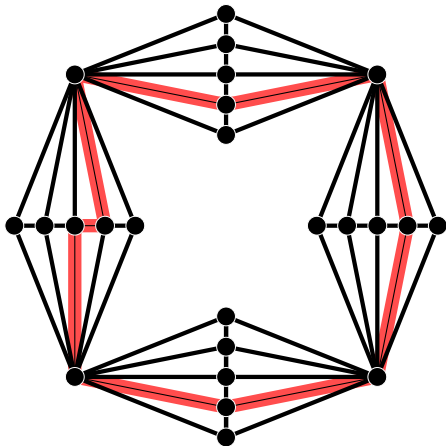
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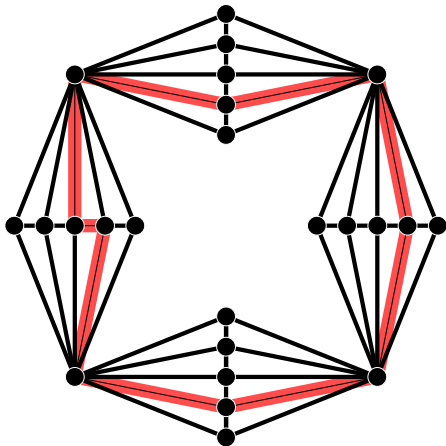
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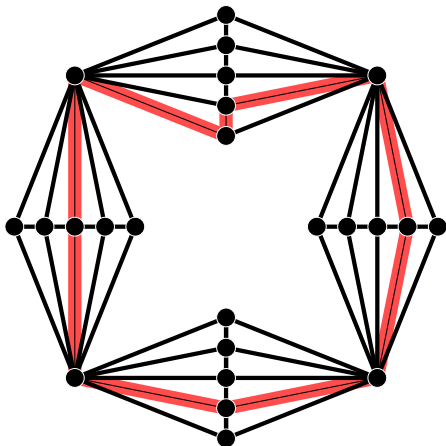
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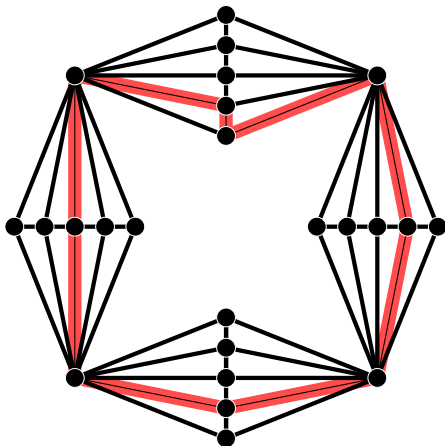
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Open Problems

Question

- *What is $N_{\mathcal{P}}(n, P_m)$ for $m = 9, 11, \dots$? [Reduction lemma exists]*
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- *What is $N_{\mathcal{P}}(n, P_m)$ for $m = 6, 8, \dots$? [No reduction lemma yet]*

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Question (General maximum likelihood)

Given a graph H on m edges.

- For a probability distribution μ on $\binom{V}{2}$, choose m edges and compute $\mathbb{P}(\text{a copy of } H) = \beta(\mu, H)$.
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Although a μ achieving the maximum often occurs, it doesn't have to.

Observation

- If $H = K_{1,m}$ then the maximum likelihood results from a sequence $\{\sigma_n\}$ of stars with each edge having probability $1/n$.
- If $H = m \times K_2$ then the maximum likelihood results from a sequence $\{\mu_n\}$ of matchings with each edge having probability $1/n$.

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Remark

- Finding $\rho(\mu; m)$, which counts odd paths is not exactly this kind of maximum likelihood question.
- The odd cycle C_{2m+1} problem requires computing $2m \cdot \beta(\mu, C_m) + \beta(\mu, P_{m+1})$.

Thank you!